CPSC 320 Notes: DP in 2-D

October 31, 2016

The Longest Common Subsequence of two strings A and B is the longest string whose letters appear in order (but not necessarily consecutively) within both A and B. For example, the LCS of eleanor and naomi is the length 2 string no (or equivalently the length 2 string ao).

(Biologists: If these were DNA base or amino acid sequences, can you imagine how this might be a useful problem?)

1. Write at least three trivial or small examples and their solutions.
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2. Consider the two strings tycoon and country. Describe the relationship of the length of their LCS with the length of the LCS of tycoon and countr (the same string A and string B with its last letter removed).

3. Now consider the two strings stable and marriage. Describe the relationship of the length of their LCS with the length of the LCS of stabl and marriag (strings A and string B with their last letters removed).

4.	Given two strings A and B of length $n > 0$ and $m > 0$, break the problem of finding the length of the
	LCS LLCS(A[1n], B[1m]) down into a recurrence over smaller problems. USE and generalize
	your work in the previous problems!

LLCS(A[1n], B[1m])	=	
the	of	
	,	
	,	and

5. Given two strings A and B, if either has a length of O, what is the length of their LCS?

6. The previous two problems give a recurrence to solve LLCS. Does this recurrence repeatedly solve suproblems many times? (That is, might we want to use memoization or dynamic programming on it?) Sketch enough of the recursion tree to justify your answer.

7.	Convert	vour	recurrence	into a	memoized	solution	to the	LLCS	problem.

8. Complete the following table to find the length of the LCS of tycoon and country using your memoized solution:

	 С	О	u	n	t	r	у
_							
t							
У							
c							
О							
О							
n							

9. Go back to the table and extract the actual LCS from it. Circle each entry of the table you have to inspect in constructing the LCS. Then, use the space below to write an algorithm that extracts the actual LCS from an LLCS table.

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// Note: len(A) = n, len(B) = m, and Table is a filled-in // (n+1)x(m+1) LLCS memoization table for A and B ExplainLCS(A, B, Table):
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10.	Give a dynamic programming solution that produces the same table as the memoized solution.
11.	Analyse the efficiency of your memoized, DP, and "explain" algorithms in terms of runtime and (additional, beyond the parameters) memory use. You may assume the strings are of length n and m , where $n \leq m$ (without loss of generality).
12.	If we only want the length of the LCS of A and B with lengths n and m , where $n \leq m$, explain how we can "get away" with using only $O(n)$ memory in the dynamic programming solution.

1 Challenge

- 1. **Prove** that if two strings end in the same letter, you can ignore all but one of the "options" in your recurrence.
- 2. Give a LCS algorithm that runs in the same asymptotic runtime as the one above, uses only O(m+n) space (note that this is potentially more than the "space-efficient" version mentioned above), and returns not only the length of the LCS but the LCS itself. (Note: try this for yourself for a while, and then walk through the description of the awesome algorithm in section 6.7 if you need help.)