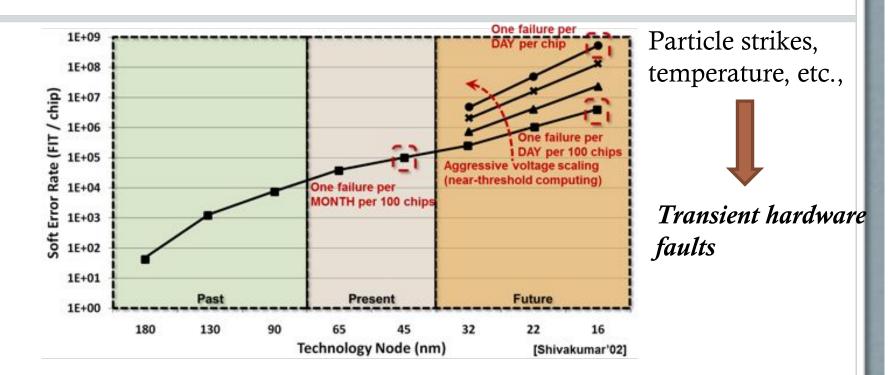
SDCTune: A Model for Predicting the SDC Proneness of an Application for Configurable Protection



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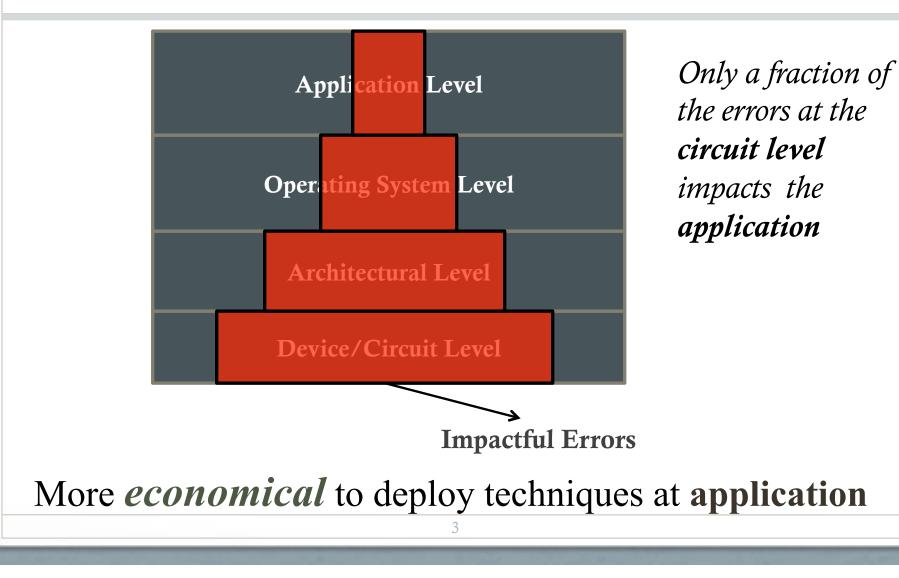
Motivation: Transient Errors

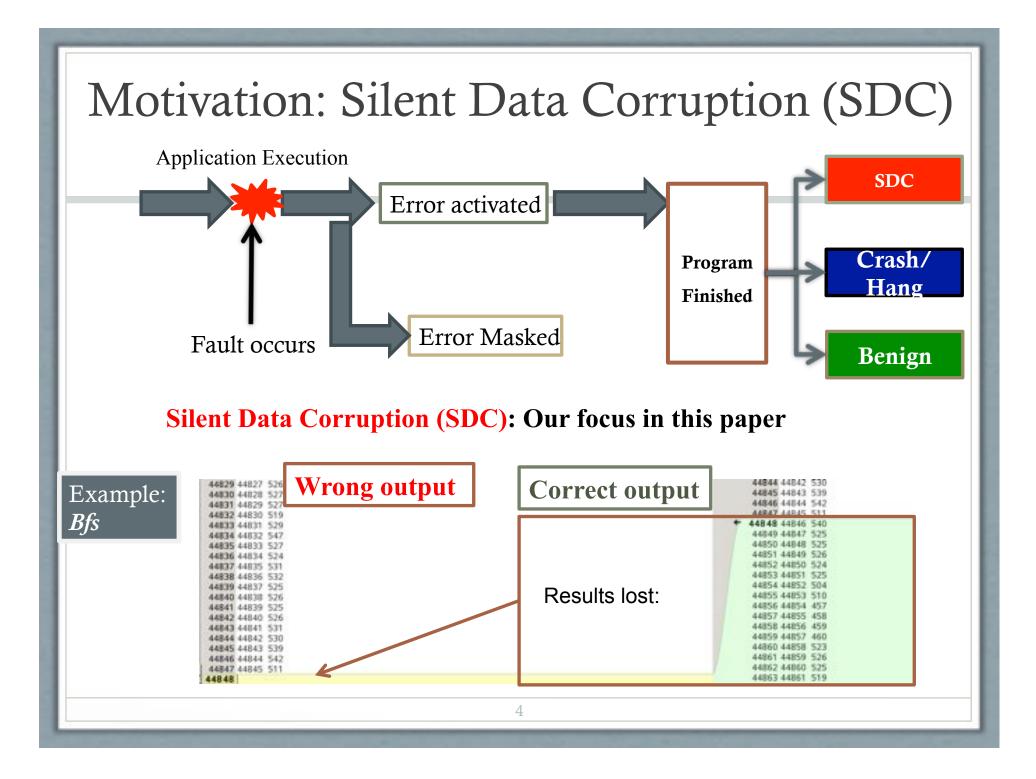


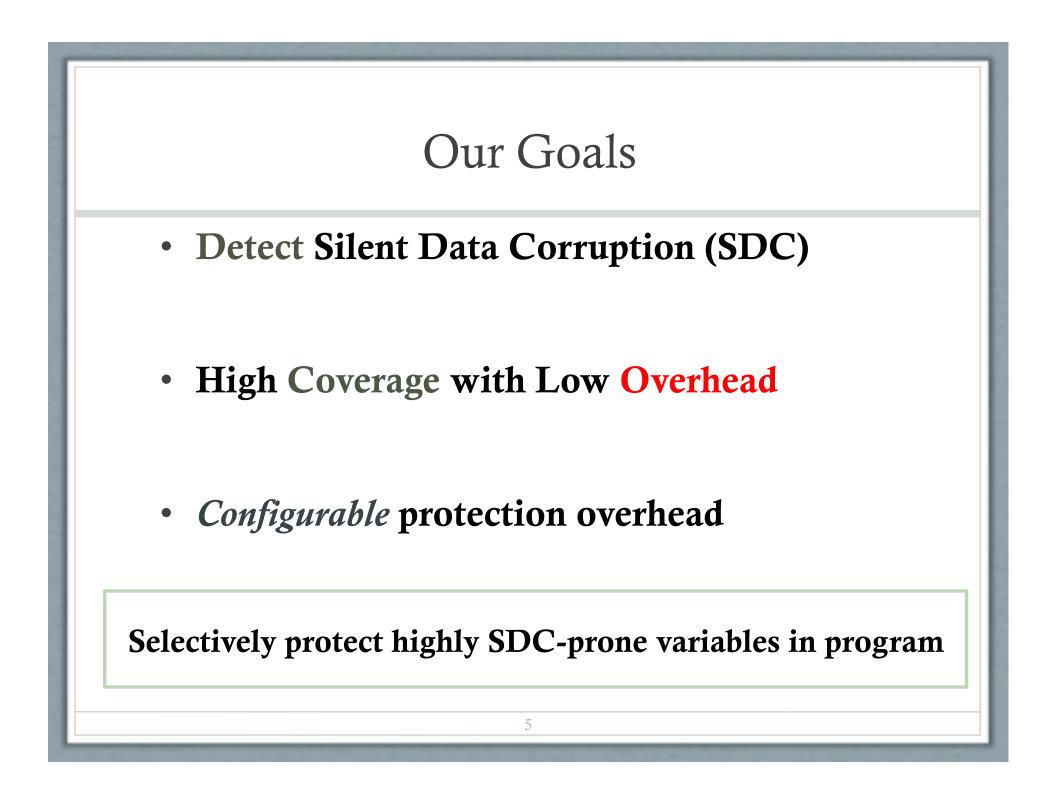
Source: Feng et. al., ASPLOS'2010

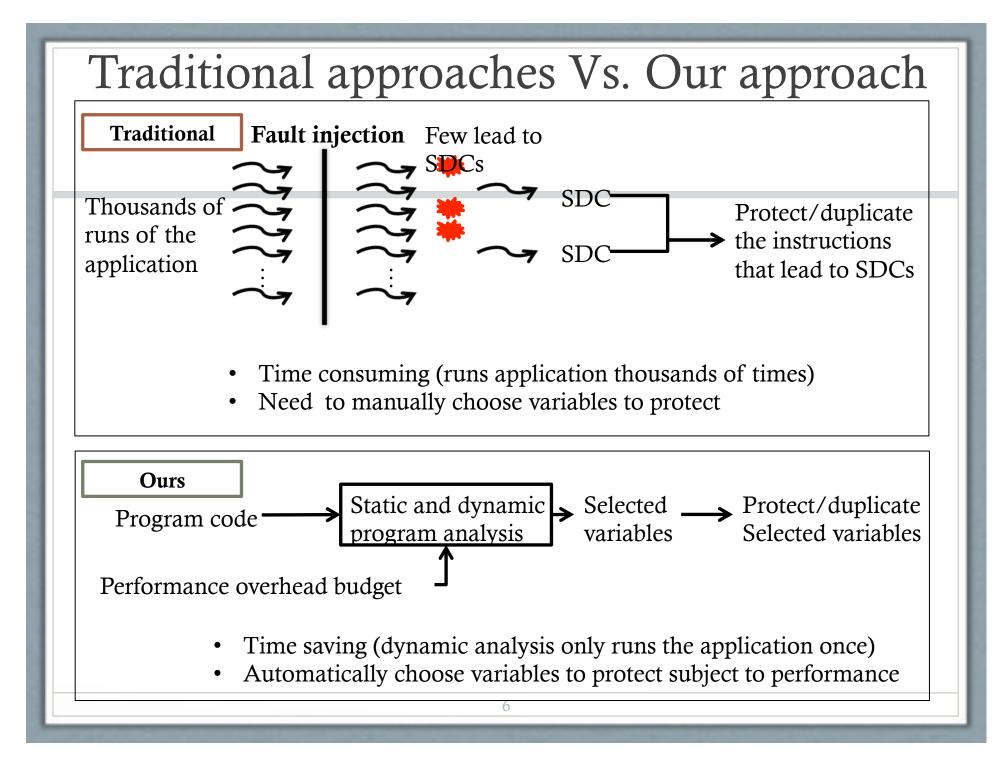
Transient hardware errors (aka. Soft errors) *increase* as **feature** sizes *shrink*

Motivation: Application-level Techniques









Fault model

• Single bit flip fault

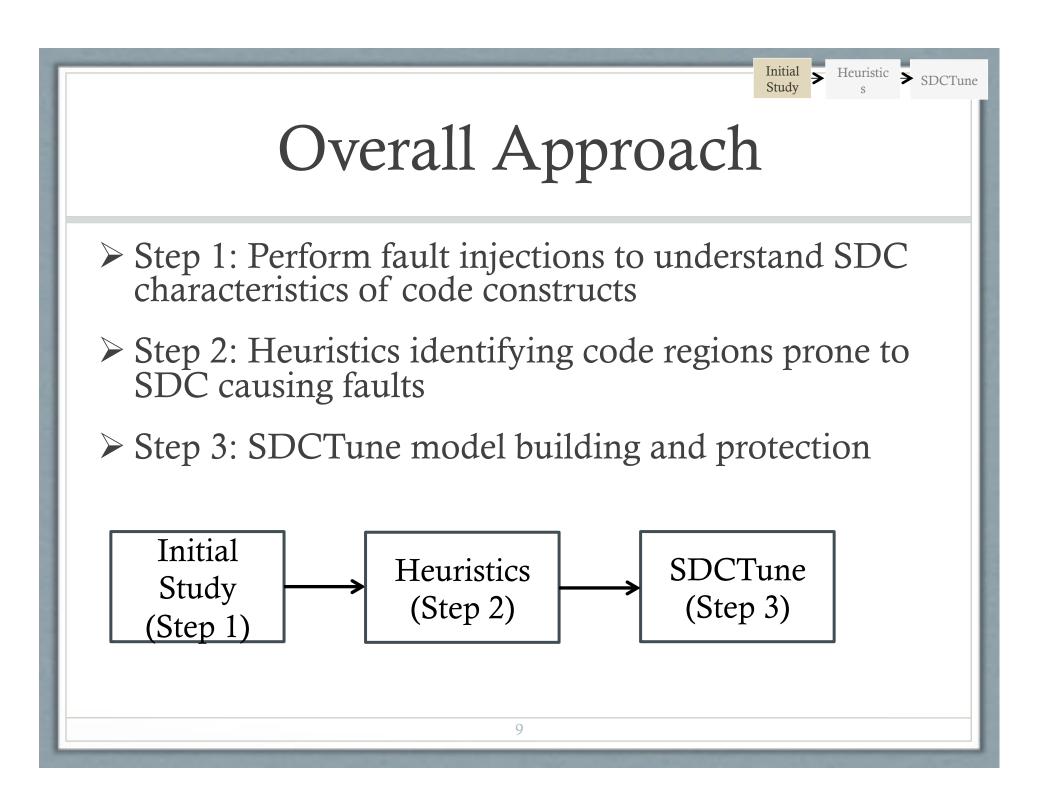
• One fault per run

• Errors in registers and execution units

• Program data that is visible at architectural level

Motivation and Goal

- Approach
- Evaluation and Results
- Conclusion



Initial study: Goals

• Initial fault injection experiments

• The goal is to understand the reasons for SDC failures

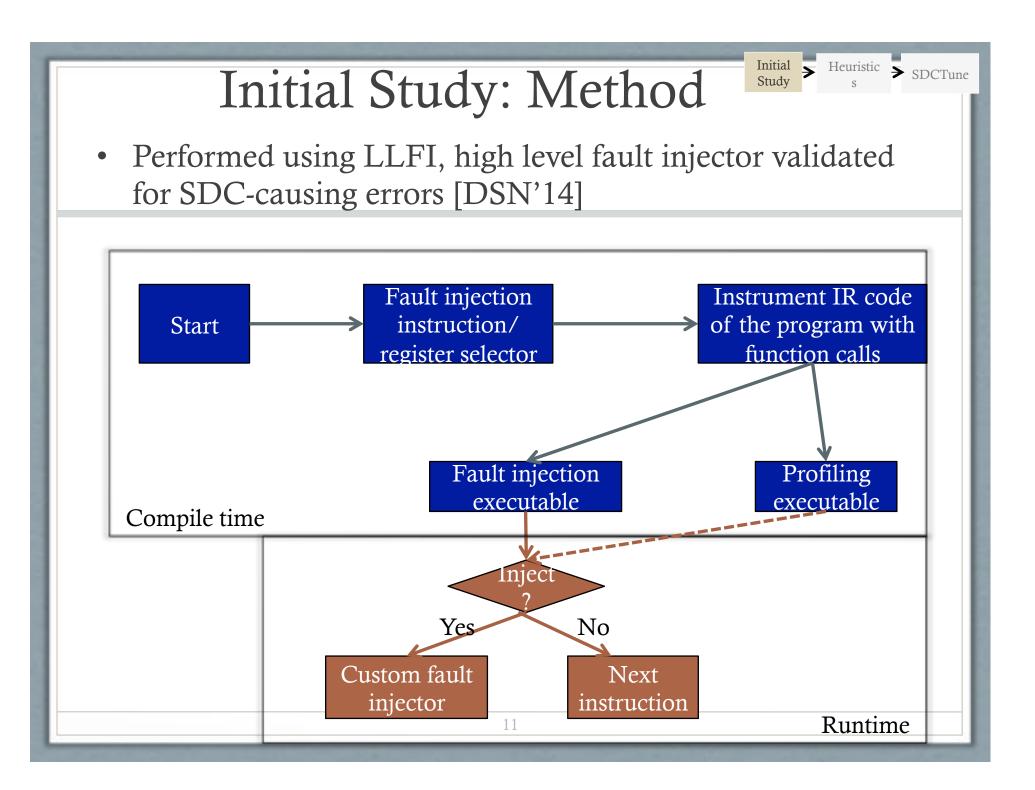
Heuristic > SDCTune

Initial Study

• Used to formulate heuristics for selective protection

Manually inspect why SDC occurs

- *Highly executed instructions cover most SDCs*
- Not all highly executed instructions should be protected
- Find common patterns used for developing heuristics



Initial study: Findings

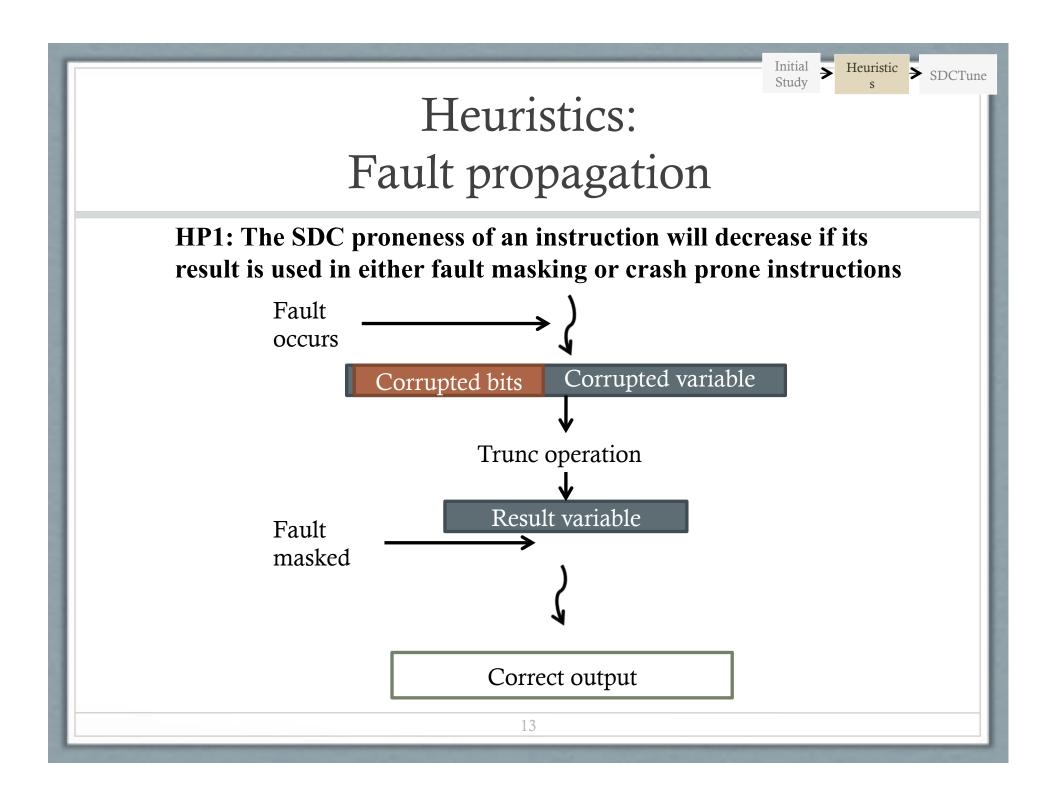
➤ Heuristic ➤ SDCTune

Initial Study

- SDC proneness of instruction depends on:
 - The fault propagation in its data dependency chain
 - The SDC proneness of the end point of that chain

- End points of data dependency chain:
 - Store operations
 - Comparison operations

Need heuristics for fault propagation, store operations, comparison operations



Heuristics: Store operations

→ Heuristic → SDCTune

Initial Study

Category	Description	Major related features Data width Data width and control flow deviation	
Addr NoCmp	The stored value is used in calculating memory addresses but not comparison results		
Addr Cmp	The stored value is used in calculating both memory addresses and comparison results		
Cmp NoAddr	The stored value is used in calculating comparison results but not memory addresses	Resilient or Unresilient comparison	
NoCmp NoAddr	The stored value is neither used in memory address calculation nor comparison results	Used in output or not	

HS1: Addr NoCmp stored values have low SDC proneness in general

HS2: Addr Cmp stored values have higher SDC proneness than Addr NoCmp </br><More heuristics in paper>

Heuristic > SDCTune Initial ⋝ Study Heuristics: **Comparison** operations HC1: Nested loop depths affect the SDC proneness of loops' comparison operations. void BZ2_hbMakeCodeLengths (...){ while(nHeap>1){ //outer loop 2 SDC proneness of "*nHeap>1*" higher than 3 while(weight[tmp]<weight[4 "weight[tmp]<weight[heap[zz>>1]] heap[zz>>1]]){ ,, // inner loop 5 Heap[zz]=heap[zz>>1]; 6 zz>>1; 8 <More heuristics in paper> 10 15

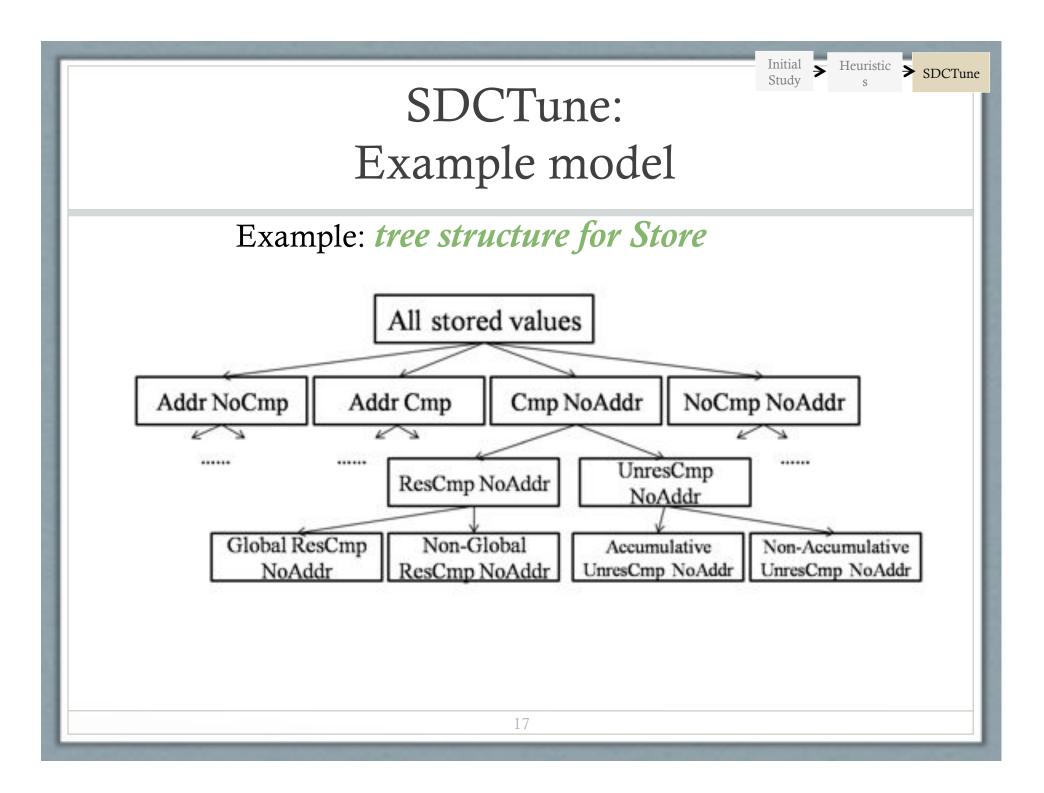
SDCTune: Build model

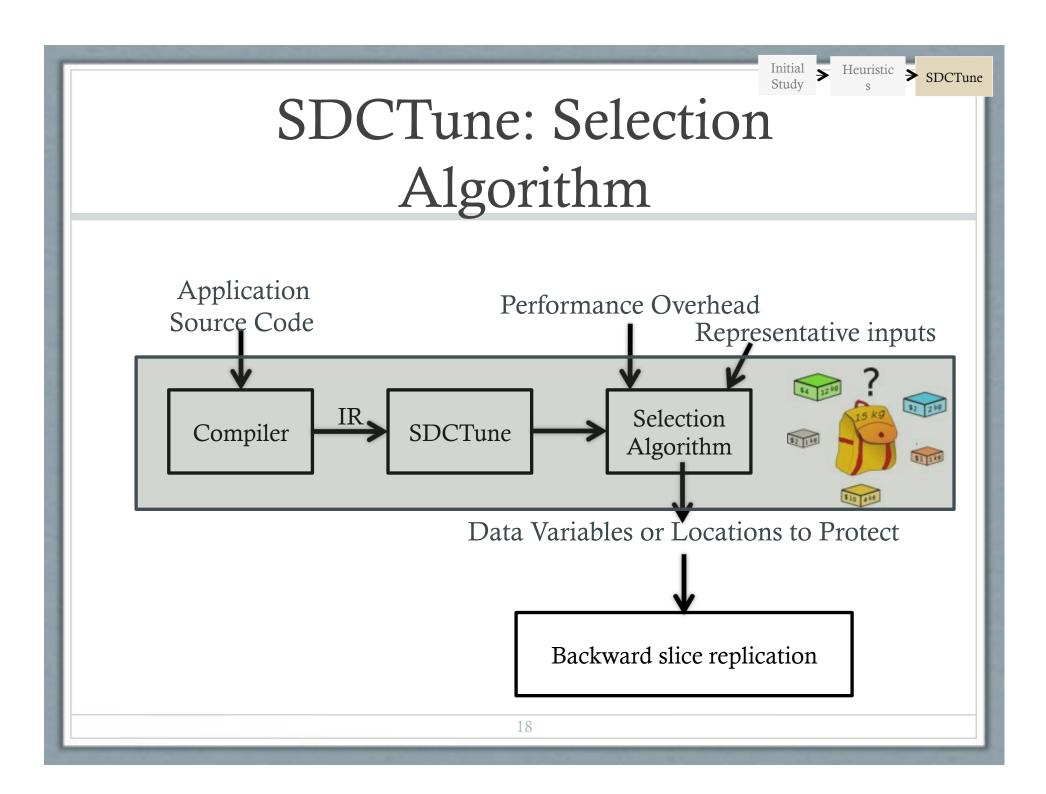
➤ Heuristic
> SDCTune

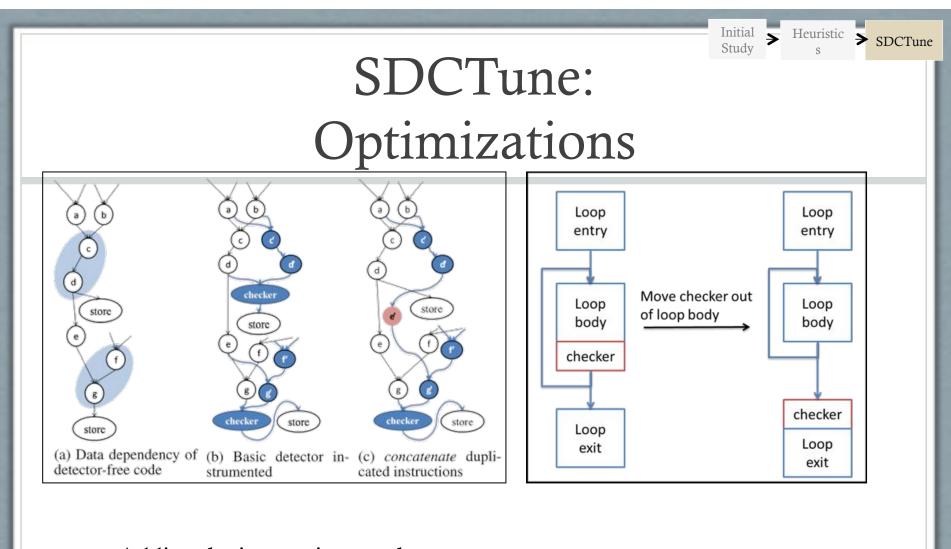
Initial Study

- Classification
 - Different types of usage are usually independent of each other
 - Classify the stored values and comparison values according to the heuristic features we observed before
- Regression
 - With same type of usages, SDC rate may show gradually correlations to several features

52 features in total used in the model





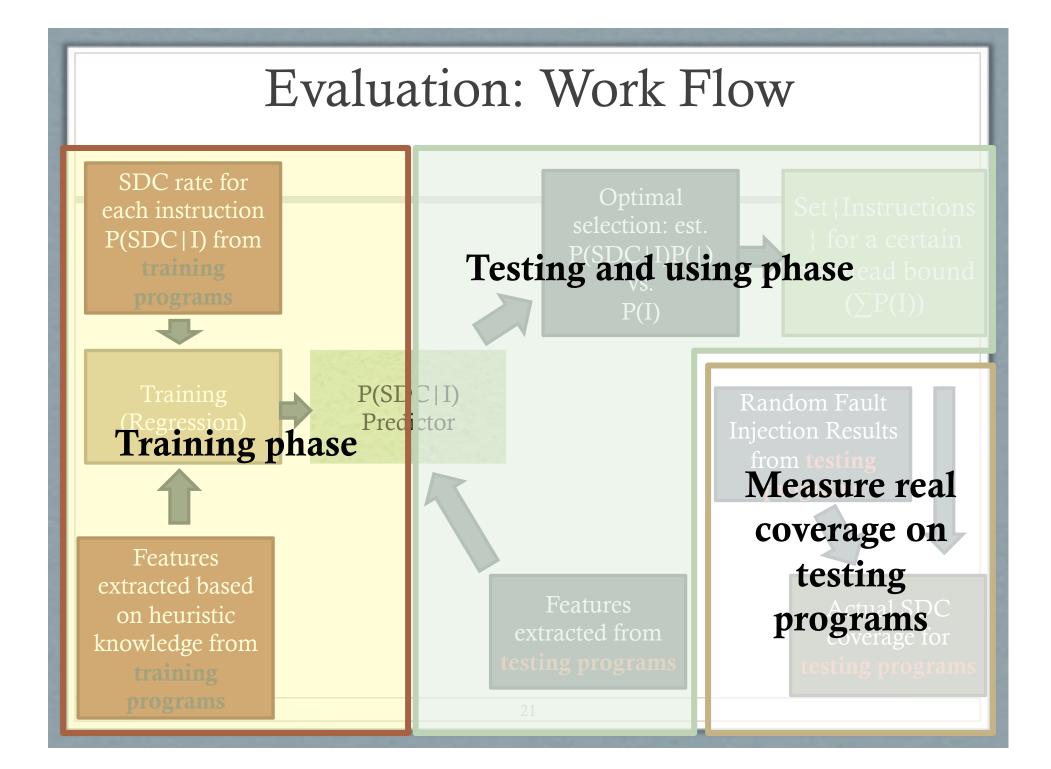


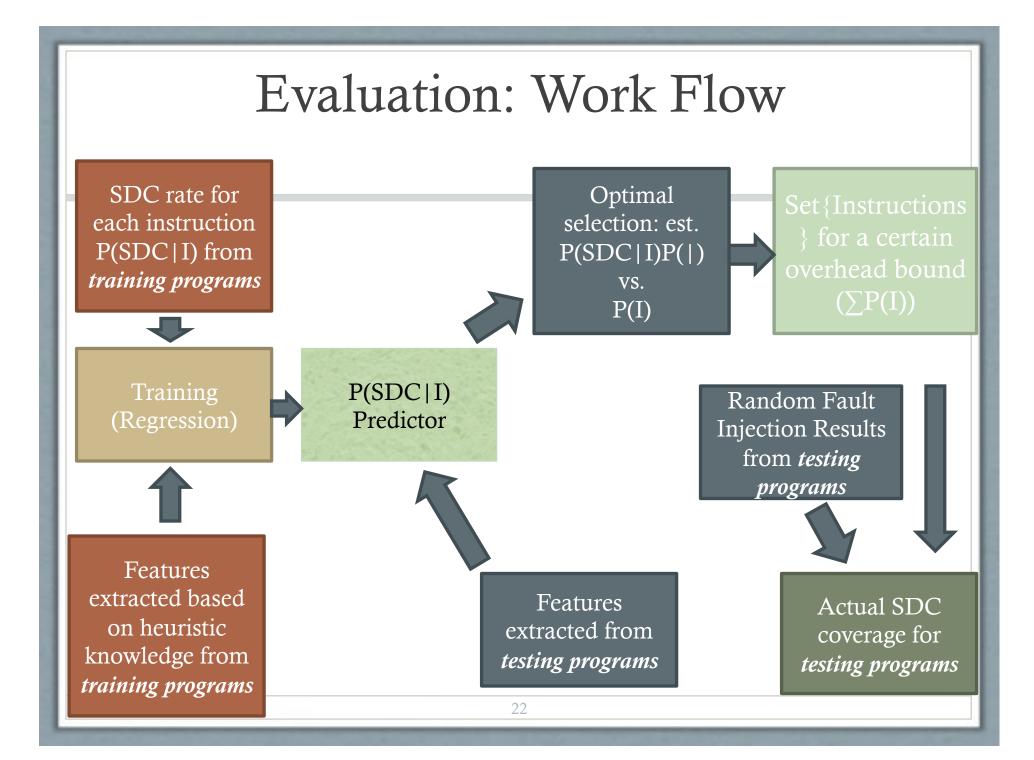
Adding the instructions to the protection set to save checkers

Move checker out of loop body

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Evaluation: Benchmarks

Training programs		Testing programs				
Program	Description	Benchmark suite	Program	Description	Benchmark suite	
IS	Integer sorting	NAS	Lbm	Fluid dynamics	Parboil	
LU	Linear algebra SPLA	SPLASH2	Gzip	Compression	SPEC	
				Large-scale		
Bzip2	Compression	SPEC	Ocean	ocean	SPLASH2	
	Price			movements		
Swaptions	portfolio of swaptions	PARSEC	Bfs	Breadth-First search	Parboil	
Water	Molecular dynamics	SPLASH2	Mcf	Combinatoria 1 optimization	SPEC	
CG	Conjugate gradient	NAS	Libquantu m	Quantum computing	SPEC	

Evaluation: Experiments

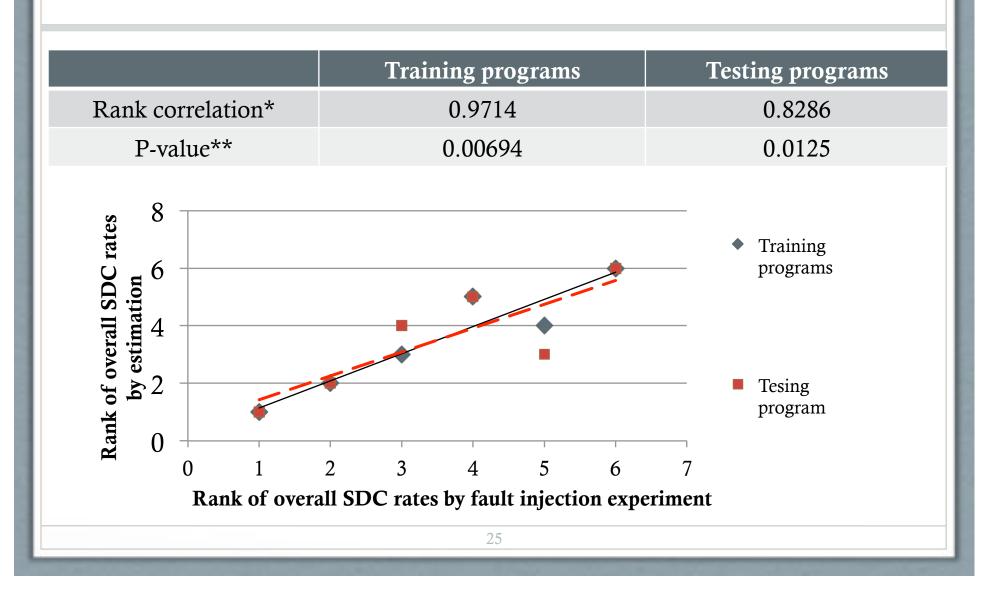
- Estimate overall SDC rates using SDCTune and compare with fault injection experiments
 - Measure correlation between predicted and actual
- Measure SDC Coverage of detectors inserted using SDCTune for different overhead bounds

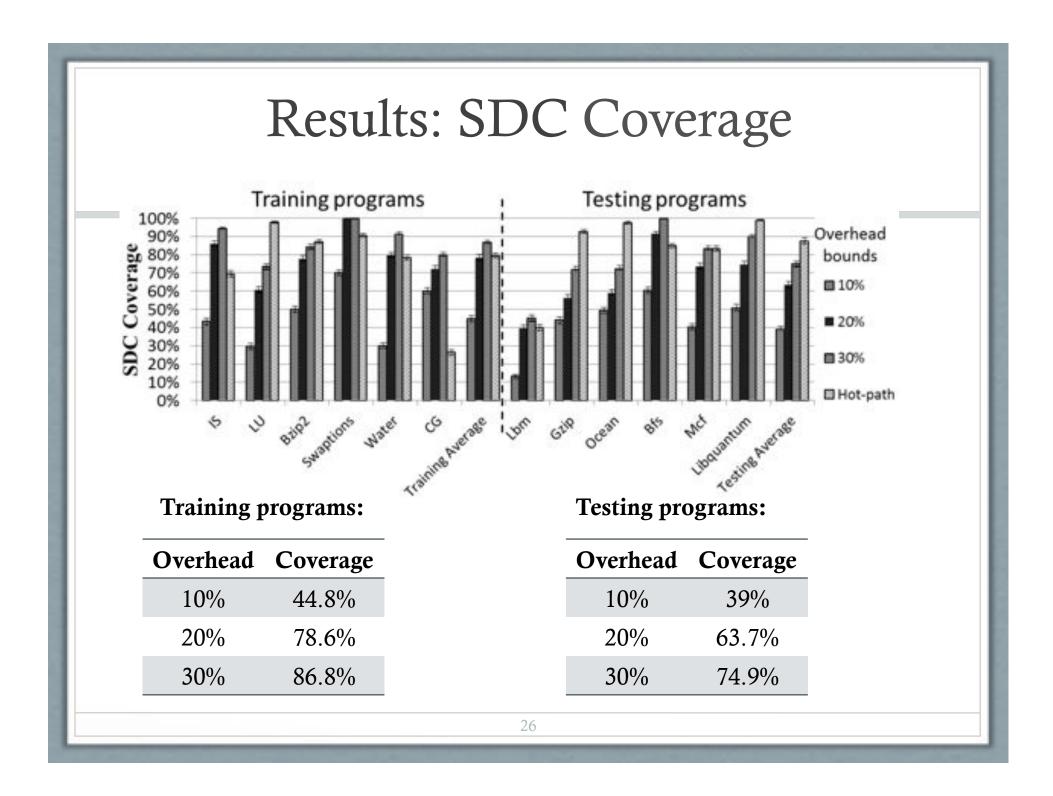
• Consider 10, 20 and 30% performance overheads

• Compared performance overhead and efficiency with full duplication and hot-path duplication

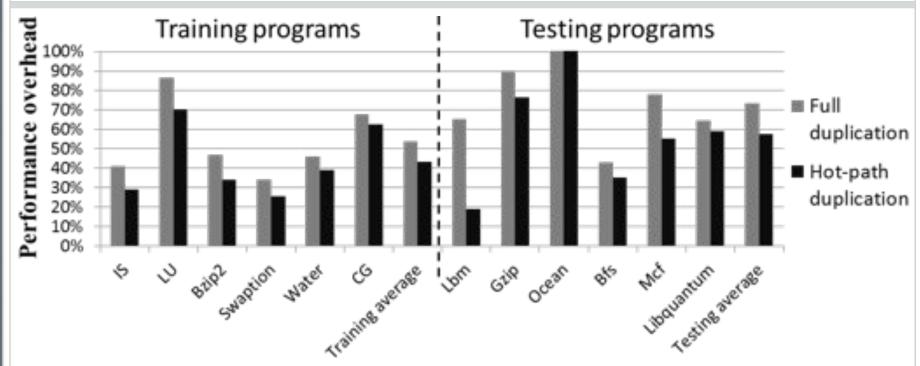
• Efficiency = SDC coverage / Performance overhead

Results: Overall SDC Rates



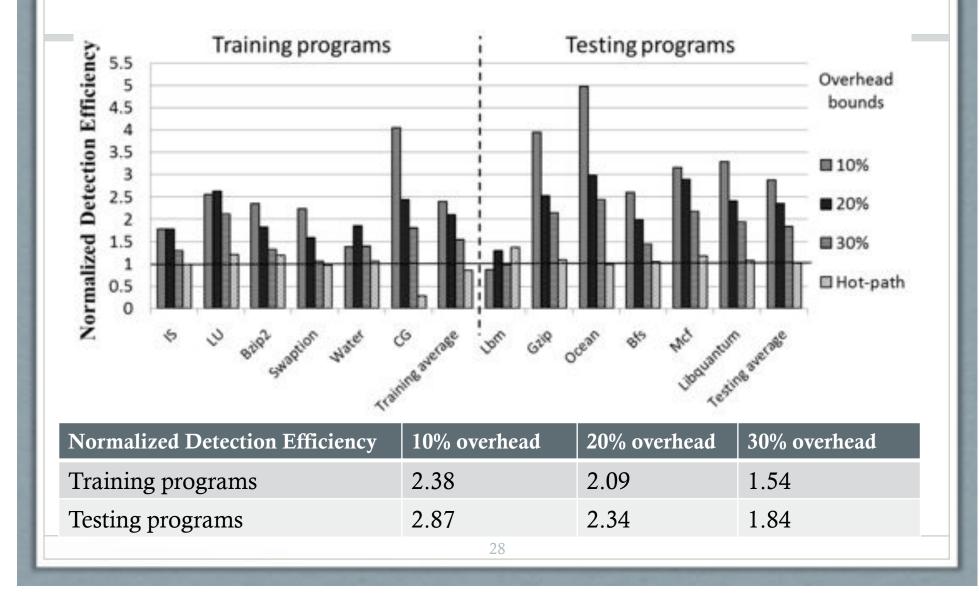


Results: Full Duplication Overheads



Full duplication and hot-path duplication (top 10% of paths) have high overheads. For full duplication it ranges from 53.7% to 73.6%, for hot-path duplication it ranges from 43.5 to 57.6%.

Results: Detection Efficiency



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Conclusion and Future Work

- Configurable protection techniques for SDC failures are required as transient fault rates increase
- We find heuristics to estimate SDC proneness for program variables **based on static and dynamic features**
- SDCTune model to guide configurable SDC protection
 - Accurate at predicting relative SDC rates of applications
 - Much better detection efficiency compared to full duplication
- Future work
 - Improving the model's accuracy using auto-tuning
 - Using symptom based detectors for protection

http://blogs.ubc.ca/karthik/