Adversary Models

EECE 571B "Computer Security"

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why we need adversary models?

attacks and countermeasures are meaningless without





elements of an adversary model

objectives

- obtain secret(s): decrypt cipher-text, guess/find password
- obtain access to assets: access to an account, full or partial control of a system or its parts

initial capabilities

- knowledge of (1) keys, passwords, and other secrets, (2) system/environment design/architecture
- access to the system's source code and other implementation details
- partial access to a system (PC, server, mobile device)
- partial control of a system (direct browser to a URL, control of a low-privilege account)

capabilities during the attack

- passive: eavesdropping messages
- active: modifying, re-playing, or removing messages
- running code on the target system
- observing system at run-time





Dolev-Yao model

- the network is completely under the adversarial control
 - can record, delete, replay, reroute, reorder, and completely control the scheduling of messages.
- the adversary is the network
 - the honest participants send their messages only to the adversary and receive messages only from the adversary.
- the adversary can choose the recipient and auxiliary information for its messages with total non-determinism
- initial knowledge of the adversary
 - the public keys (K_{Pub}),
 - the private keys of subverted participants (K_{Adv} ⊆ K_{Priv}),
 - the identifiers of the principals (I), and
 - the nonces the adversary itself generates (R_{Adv} ⊆ R), which are assumed to be distinct from all nonces generated by honest participants.





Dolev-Yao model (continued)

- message M is derivable by adversary from a set of messages S, if it's possible to produce by applying the following operations a finite number of times:
 - decryption with known or learned private keys
 - encryption with public keys
 - pairing of two known elements
 - separation of a pair into its components





Chip & PIN







EMV protocol

Europay, MasterCard, VISA (EMV) -- protocol for payment cards with chips (and PINs)

750M cards currently deployed

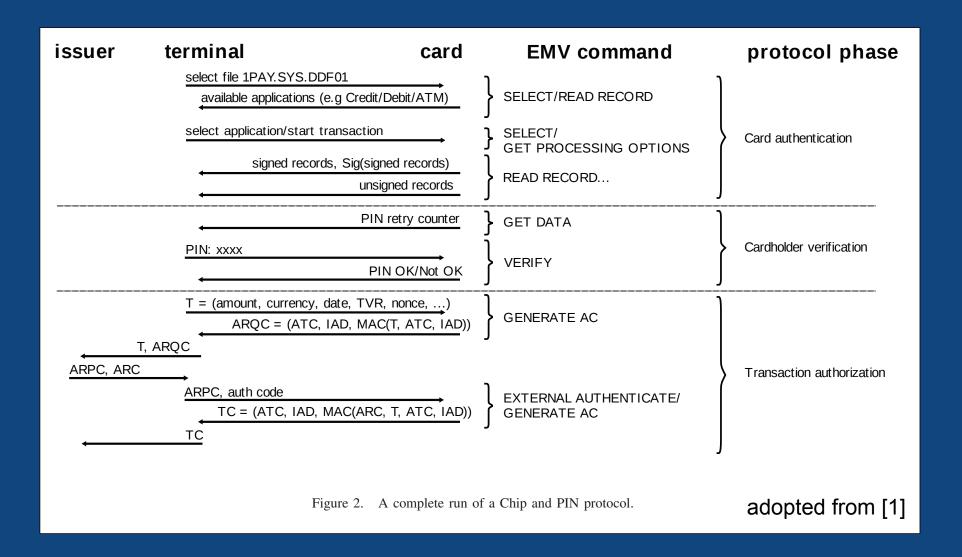
a three phase protocol:

- 1. Card authentication
 - type of card, issuer, verification method list etc)
- 2. Cardholder verification, based on verification method list,
 - 1.PIN
 - 2.signature
 - 3.nothing
- 3. Transaction authorization
 - card generates secured transaction info for the issuing bank clearance





a complete run of a Chip & PIN protocol







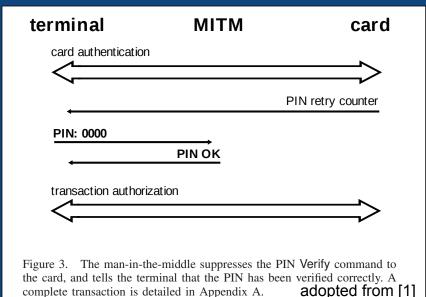
video clip

http://www.youtube.com/watch?v=JPAX32lgkrw





cardholder verification step



 attacker tricks the card into "thinking" it's doing a chip-andsignature transaction while the terminal "thinks" it's chipand-PIN.





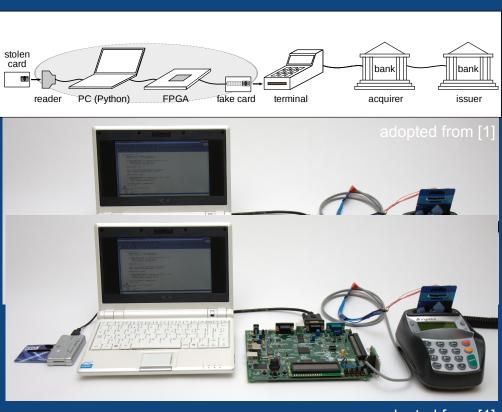


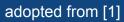
adopted from [1]





the attack









adversary model

objectives

- pay to a street merchant with a stolen payment C&P card
- initial capabilities
 - can still payment C&P cards
 - can purchase or make necessary equipment for the MITM attack
- capabilities during the attack
 - conceal the equipment from the merchant's staff
 - conceal the fact that the fake card has wires attached to it
 - insert the fake card in the merchant's terminal





Sony CD-DRM Episode





who is the adversary?

- user who wants to rip a CD
- researchers, who analyzed the Sony CD DRM
- Label vendor, who wants to use DRM for free
- DRM vendor, who wants to install its software on the user's PC
- criminals who want to use attack user's PC using vulnerabilities in the DRM software





Sony CD DRM: options for attack tactics

- 1. defeat the passive protection
- 2. stop the DRM software from installing itself
- 3. trick the recognition algorithm
- 4. defeat the active protection software's blocking
- 5. capture the music from the DRM vendor's player
- 6. uninstall the protection software
- 7. create state snapshot, rip CD, rollback the state





Sony CD DRM: attack adversary model

objectives

- wants to copy the music illegally
- wants to make uses allowed by copyright law but blocked by the DRM
- initial capabilities (not all required)
 - approximately knows CD structure (music and data "sections")
 - can prevent parts of the CD from being read by the drive (masking tape or felt-tip marker)
 - can configure their PC not to do autorun
 - can rename the executable of the ripping/copying application
 - has access to protected and unprotected versions of the same album and can perform binary comparison





Sony CD DRM: adversary model capabilities during the attack (not all required)

- physical access to the computer
 - can hold the keyboard Shift Key when inserting
- can play the CD in a different OS (Linux or MacOs)
- can repeatedly insert and ejecting the CD many times
- can "kill" the installer process (using the Windows Task Manager) before it can eject the CD
- can issue commands directly to the CD drive, bypassing Windows CD driver
- can log all CD device activities
- can create new watermarked disks
- can convert the tracks to a lossy format, like MP3, and then burn them back to a CD
- can flip the least significant bit of one carefully chosen sample from each of the 30 watermark clusters
- can transplant the 3 least significant bits of each sample within the watermarked region of a protected track to the corresponding sample from an unprotected track
- can deduce the full details of the structure of the watermark
- can create and restore and snapshot of the PC state
- can modify %windir%\system32\ \$sys\$filesystem\\$sys\$parking file





Question to ponder

When you read the Sony DRM Episode paper:

• What's the model of the adversary that is trying to compromise user's computer by exploiting vulnerabilities due to the DRM?



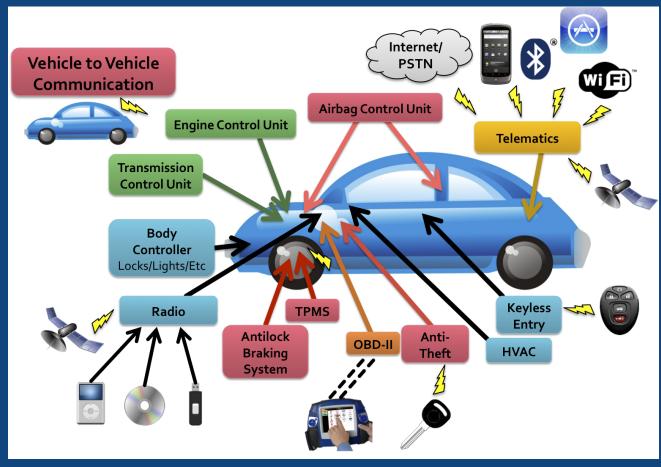


Security Analysis of a Modern Car





today cars



adopted from [2]





indirect physical access: media player attack

- attack 1: vestigial radio reflash from CD code
- attack 2: WMA parsing bug -> buffer overflow
- on-radio debugger
- insert CD containing malicious WMA file
- compromise the car





short-range wireless: Bluetooth attack

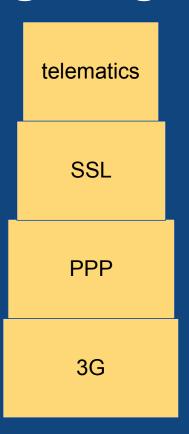
- common embedded Bluetooth stack on telematics unit
 - strcpy() bug
- Android trojan compromises telematics ECU
- can undetectably pair a bluetooth device
 - USRP-based software radio
 - brute force PIN
 - cannot be unpaired with standard interface

L





long-range wireless: cellular attack



telematics

buffer overflow

software modem

voice channel

cell phone

- call telematics unit
- transmit malicious payload (using modem protocol or just play malicious sound track over phone)





what's next?

- remotely trigger code from prior compromise
 - proximity trigger
 - broadcast trigger (FM RDS)
 - short-range targeted trigger (Bluetooth)
 - global targeted trigger (cellular)





what can an adversary do with this?

car theft

- 1. compromise car
- 2. locate it via GPS
- 3. unlock doors
- 4. start engine
- 5. bypass anti-theft
- video demo: http://www.youtube.com/watch?v=bHfOzilwXic (minute #16)

surveillance

- 1. compromise car
- 2. continuously report GPS coordinates
- 3. stream audio recorded from the in-cabin mic





adversary model

objectives

• take control over parts or the whole car in order to perform surveillance, theft, or cause car accident.

initial capabilities

- access to equipment and documentation to develop and test an attack
 - extract device's firmware
 - reverse engineer firmware
 - identify and test vulnerable code paths
 - weaponize exploits

capabilities during the attack (one of the three)

- indirect physical access to the car
 - interacts with a physical object that interacts with the car
 - diagnostic tool that plugs directly into OBD-II port
 - entertainment systems (CD player, digital multimedia port, iPod Out)
- short-range wireless signals (between 5 and 300 meters)
 - Bluetooth, Remote Key Entry, RFID car keys, Tire Pressure Monitoring Systems, WiFi, Dedicated Short Range Communications
- long-range wireless signals (greater than 1 km)
 - broadcast channels: GPS, satellite radio, digital radio, Radio Data System, Traffic Message Channel
 - addressable channels: remote telematics systems





summary: adversary model

- objectives
- initial capabilities
- capabilities during the attack





references

- 1. Chip and PIN is Broken, Murdoch, Steven J.; Drimer, Saar; Anderson, Ross; Bond, Mike; , "Chip and PIN is Broken," 2010 IEEE Symposium on Security and Privacy (SP), pp.433-446, 16-19 May 2010, doi: 10.1109/SP. 2010.33
- 2. "Comprehensive Experimental Analyses of Automotive Attack Surfaces," S. Checkoway, D. McCoy, B. Kantor, D. Anderson, H. Shacham, S. Savage, K. Koscher, A. Czeskis, F. Roesner, T. Kohno, USENIX Security, August 10–12, 2011.



