

Quantifying the Accuracy of High-Level Fault Injection Techniques for Hardware Faults

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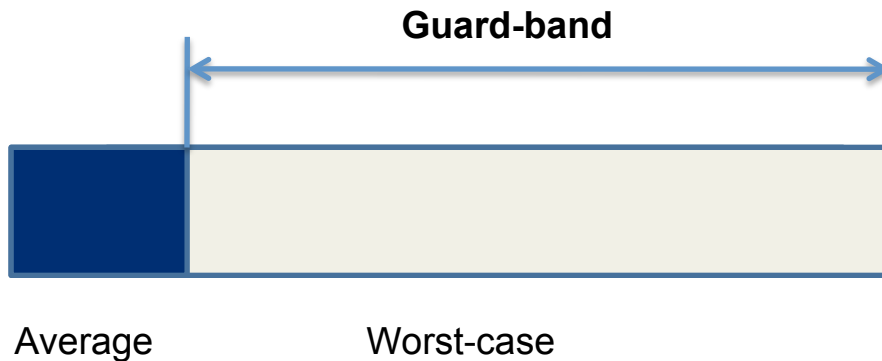
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Hardware Errors: Traditional “Solutions”

- **Guard-banding**

Guard-banding wastes power as gap between average and worst-case widens due to variations



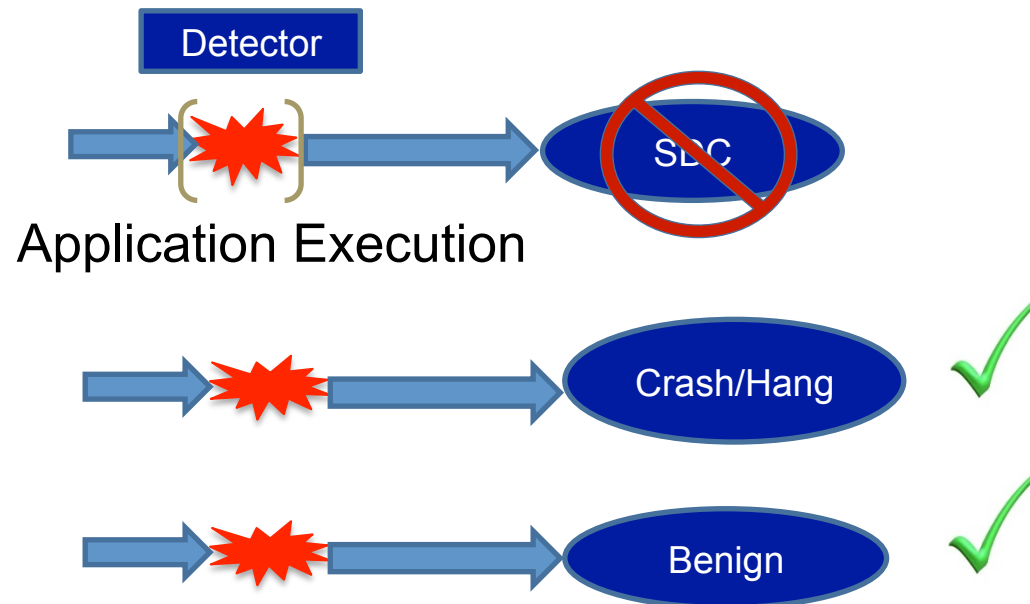
- **Duplication**

Hardware duplication (DMR) can result in 2X slowdown and/or energy consumption



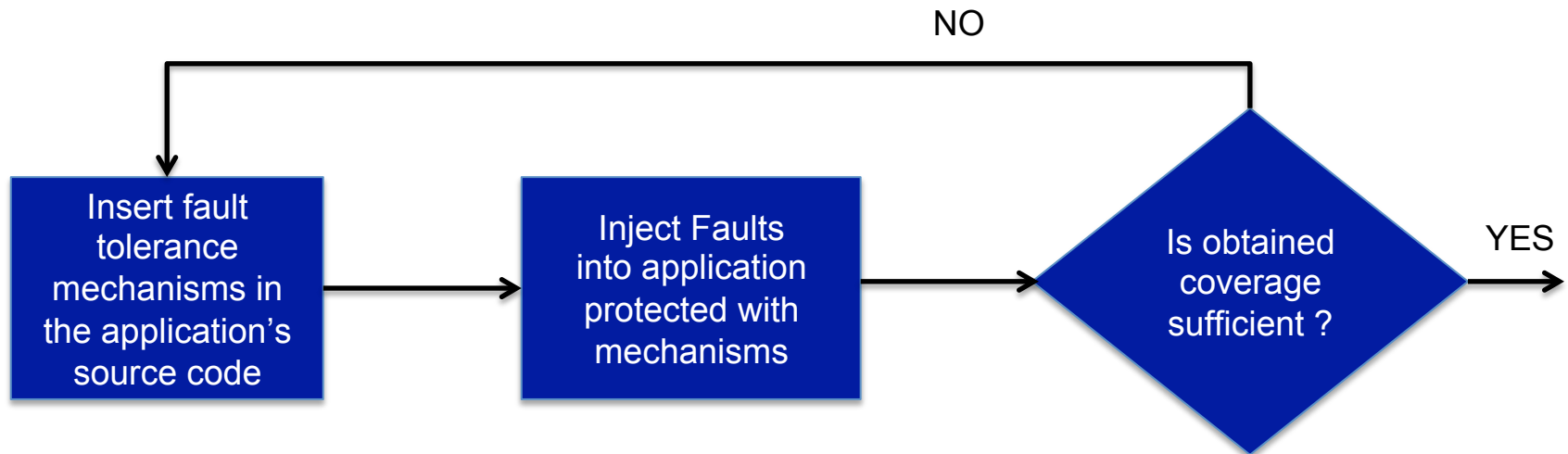
Our Research: Application-level Selective Fault-Tolerance

- Add detectors to applications to selectively detect errors causing Silent Data Corruption (SDCs) i.e., incorrect outputs



Application-level Fault Injection

- To obtain coverage estimates for applications
- Iteratively improve coverage based on the errors missed by fault tolerance mechanisms
 - Analyze the errors that are missed by the FTMs



Low-level Fault Injection

- **Inject faults into programs at the assembly code level e.g., NFTAPE, FERRARI, GOOFI, Xception**
- **Pros:**
 - Accurate at emulating hardware faults in registers, instructions and computation units (e.g., ALUs)
- **Cons:**
 - Difficult to map injection results back to source code
 - Difficult to inject faults into selected source data

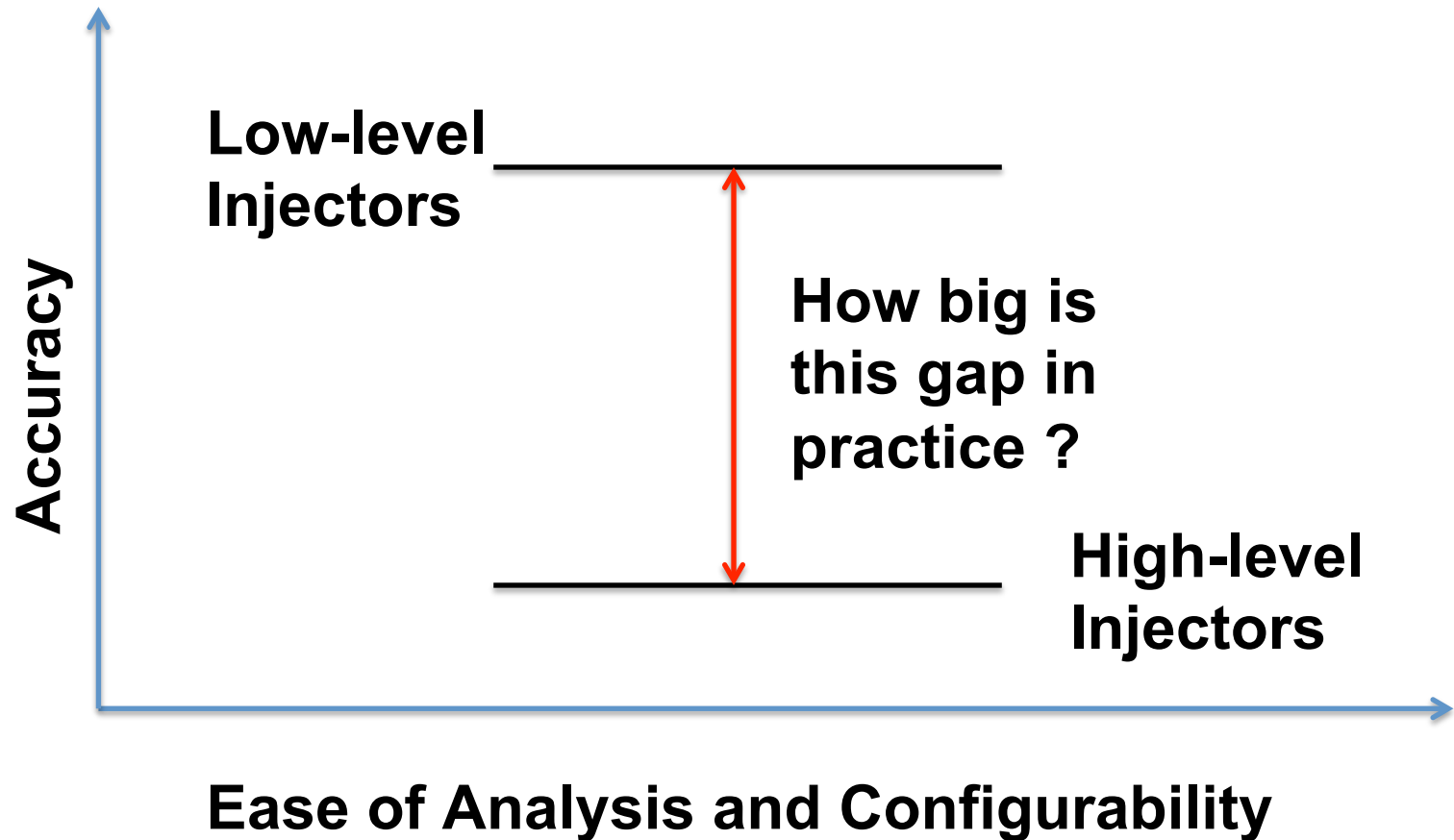
High-Level Fault Injection

- **Inject faults directly at the source code or similar levels e.g., PROPANE, Relax, Kulfi**
- **Pros:**
 - Easy to map back injection results to source code
 - Ability to inject faults into specific data-types
- **Cons:**
 - Difficult to emulate hardware faults accurately

High-Level Fault Injection: Reasons for Potential Inaccuracies

- **Lack of one-to-one mapping**
 - A single source code statement may map to multiple assembly code statements (e.g., pointers)
 - Some source statements have no analogue in the assembly code (e.g., type-cast statements)
- **Hidden States**
 - Many elements in assembly code cannot be seen in the source code (e.g., stack manipulation code)

High-Level Vs. Low-Level Injectors: Accuracy Comparison



Related Work

- **Software Faults [Madeira00][Natella13]**
 - Emulate software faults at the assembly code level
 - Inverse of our problem, as software faults occur in the source code level and are more accurate at that level
- **Safety-critical systems error consequences [Skarin-EDCC08][Pattabiraman-DSN08]**
 - Examine consequences of not considering faults at the assembly language level in design of FT mechanisms
 - Do not quantitatively measure how much the gap is

This Paper: Research Question

- How **accurate** is fault injection at the **high-level** (i.e., source code or similar levels) compared to fault injection at the **low-level** (i.e., assembly code or similar levels) ?
 - For different kinds of failures (e.g., crashes, SDCs)
 - For different kinds of instructions (e.g., loads)

Our Approach

- **Build a high-level fault injector to inject faults at the LLVM compiler's IR level: LLFI**
- **Build a low-level fault injector to inject faults using Intel's PIN tool: PINFI**
- **Compare the outcomes of LLFI and PINFI by injecting similar faults into benchmarks**

Fault Model

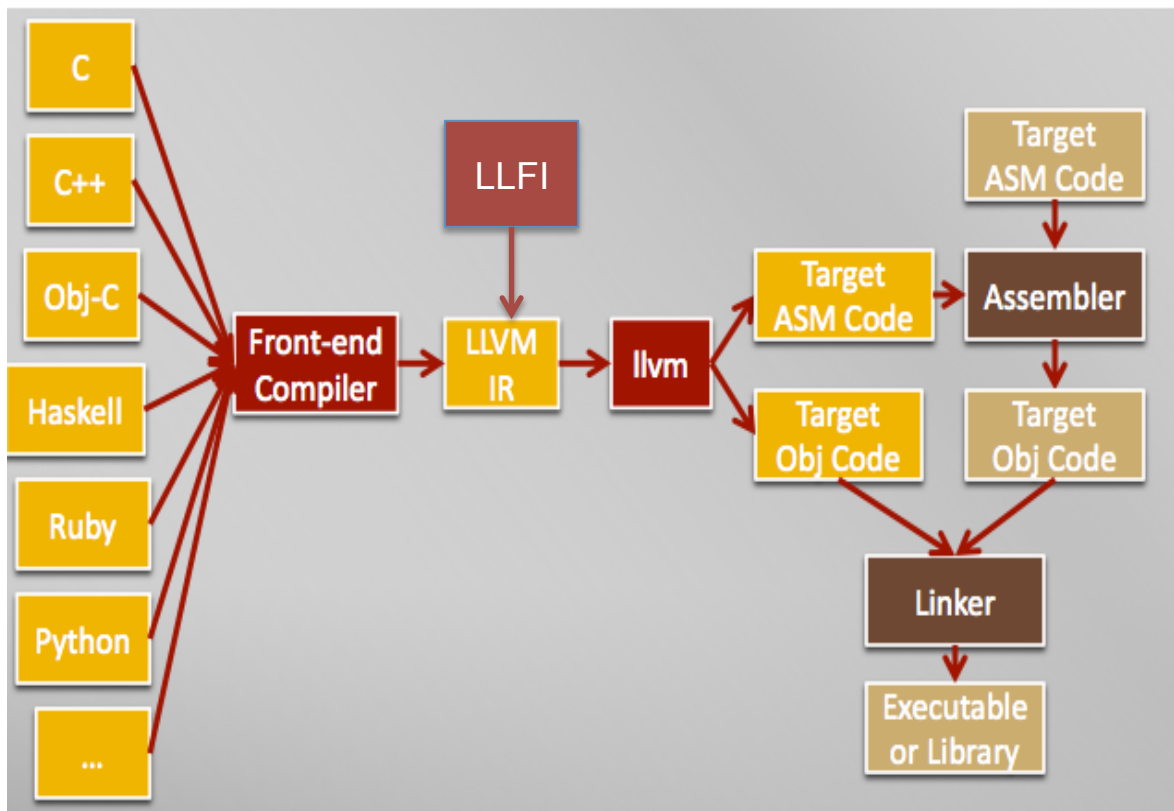
- Single bit-flip in the destination registers of a single dynamic instruction in the program
- Models transient faults in the computational parts of the processor (e.g., ALU, registers)
- Does not model memory/cache faults – assumes that these are ECC-protected
- Does not model faults in the instruction encoding

Outline

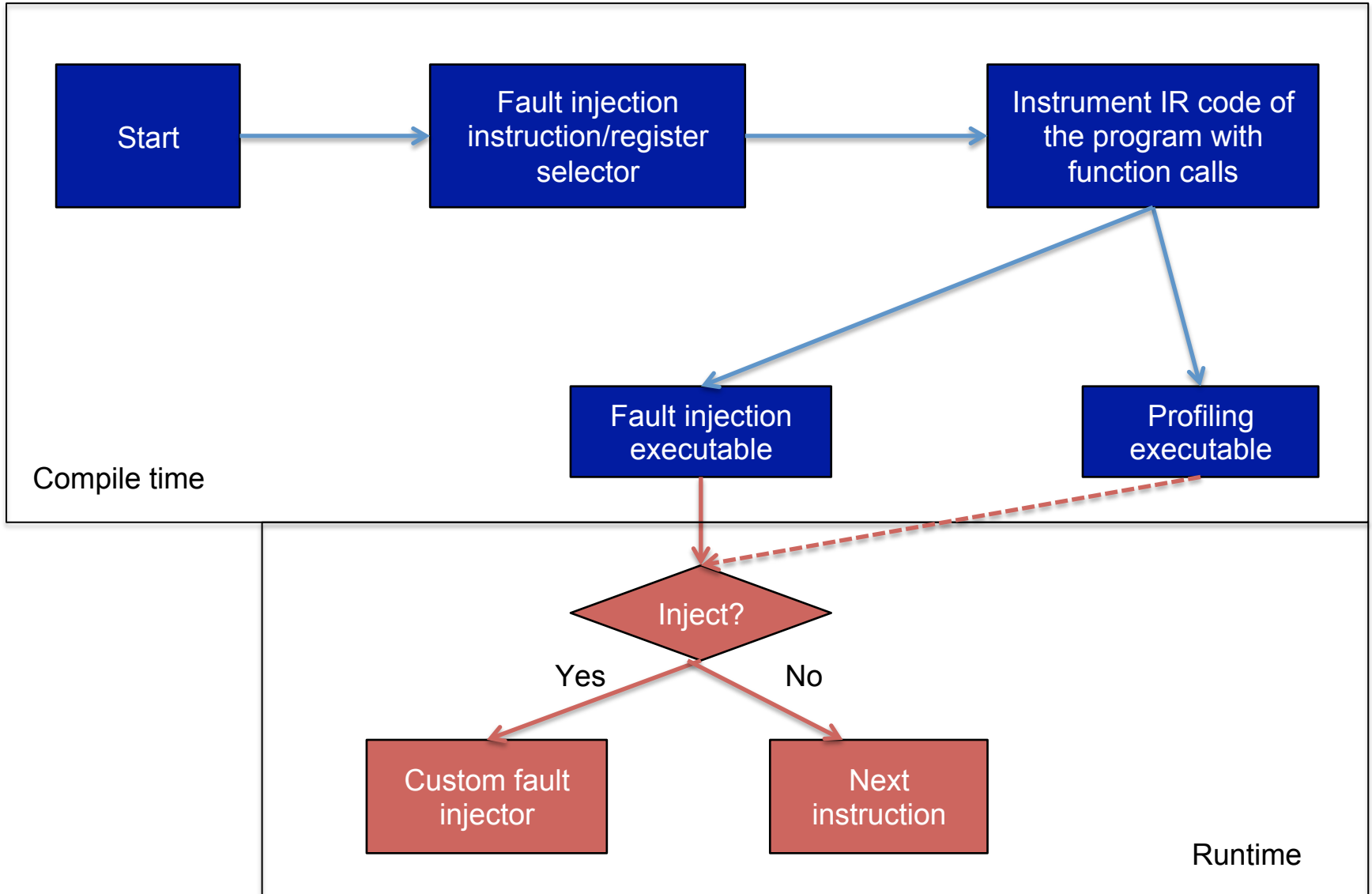
- Motivation and Approach
- **LLFI Architecture and Operation**
- PINFI Architecture and Operation
- Experimental Evaluation
- Conclusions

LLVM Fault Injector: LLFI

Works at LLVM compiler's intermediate (IR) code level [Lattner'05] – LLVM widely used in industry



How does LLFI work ?

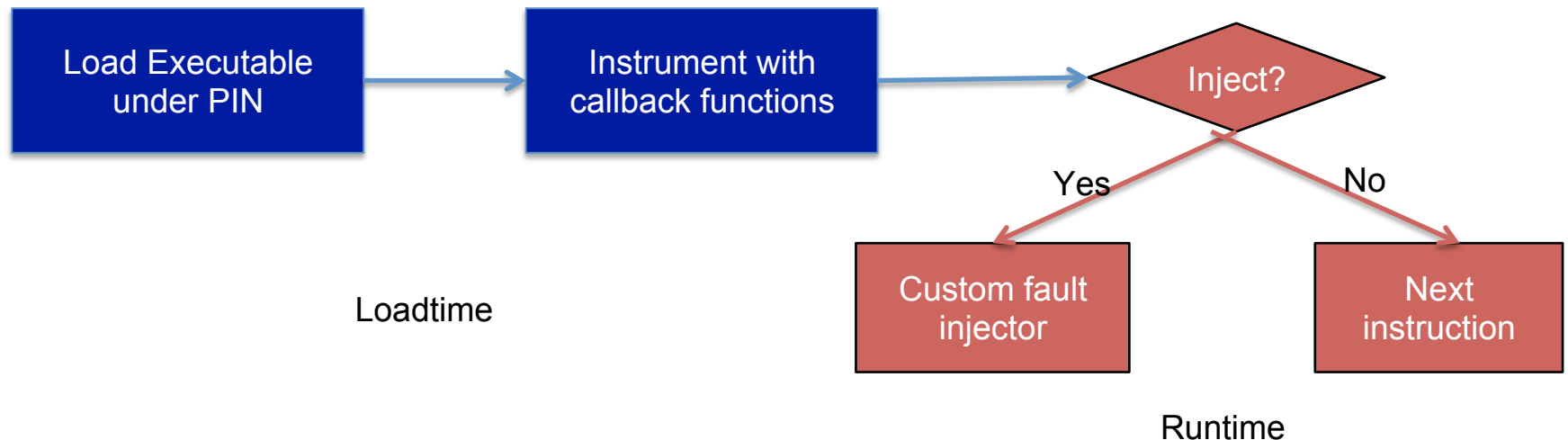


Outline

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PINFI Architecture

- Built using Intel's PIN tool for dynamic binary analysis [Luk-2005]
- Modifies executable to inject faults at runtime



Corner Cases in x86 Assembly

- Branch conditions: Flags Register

LLVM IR	X86 Assembly
<code>%11 = icmp sle i32 %9, %10 br i1 %11, label %bb, label %bb2</code>	<code>cmp \$0xa4, %eax //sets %rflags jl 4006e0</code>

- Floating point operations: Data Width

LLVM IR	X86 Assembly
<code>%3 = fadd double %1, %2</code>	<code>addsd %xmm2, %xmm0</code>

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Experimental Setup

- **Fault Injection**

- Single bit-flip in the result of a dynamic instruction
- 1000 injections per benchmark, per instruction category

- **Benchmarks**

- Four SPEC2006: bzip2, libquantum, hmmer, mcf
- Two SPLASH-2: ocean, raytrace

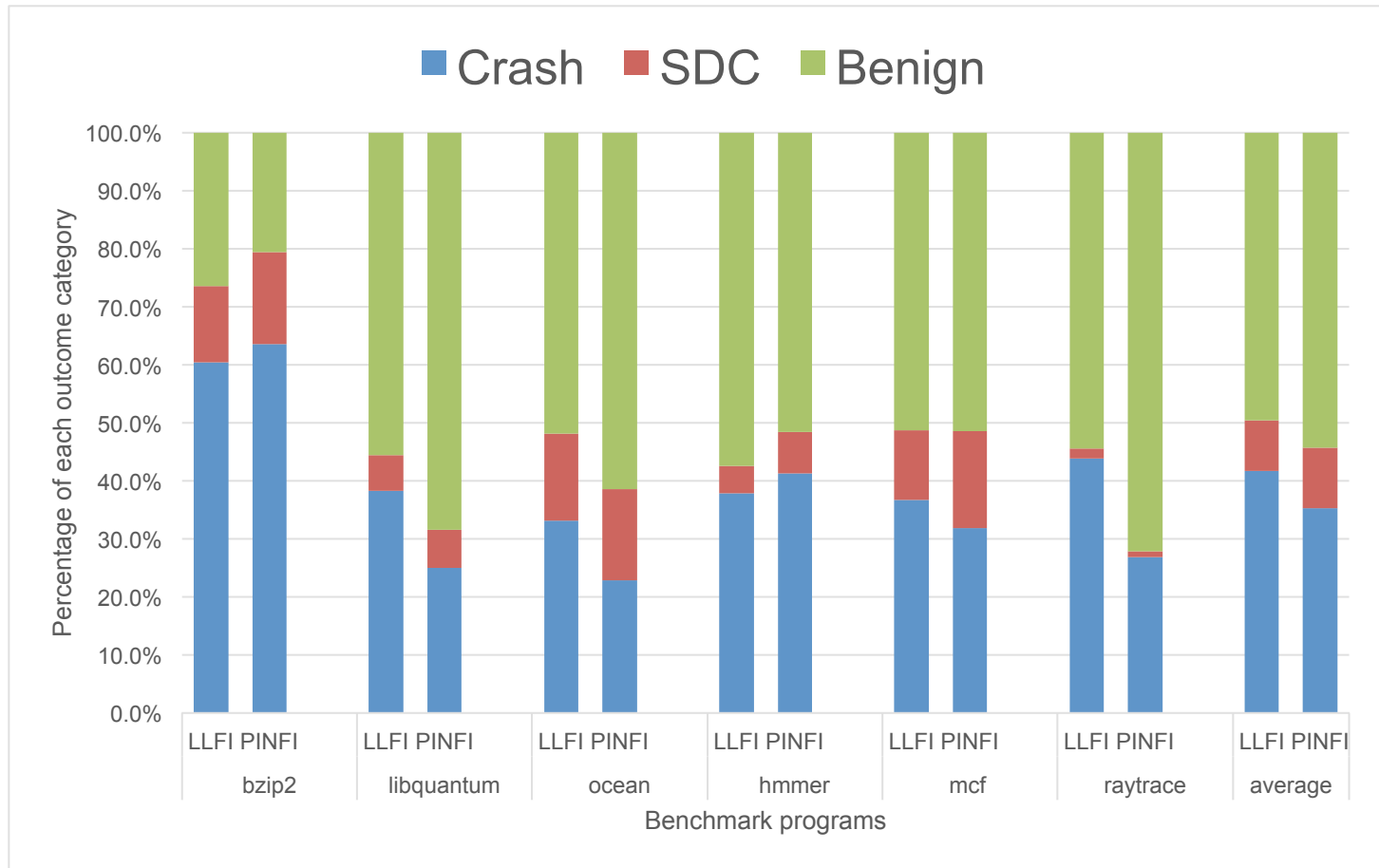
- **Outcomes**

- Crash, Hang, Benign and Silent data corruption (SDC)
- SDCs measured by comparing with golden output

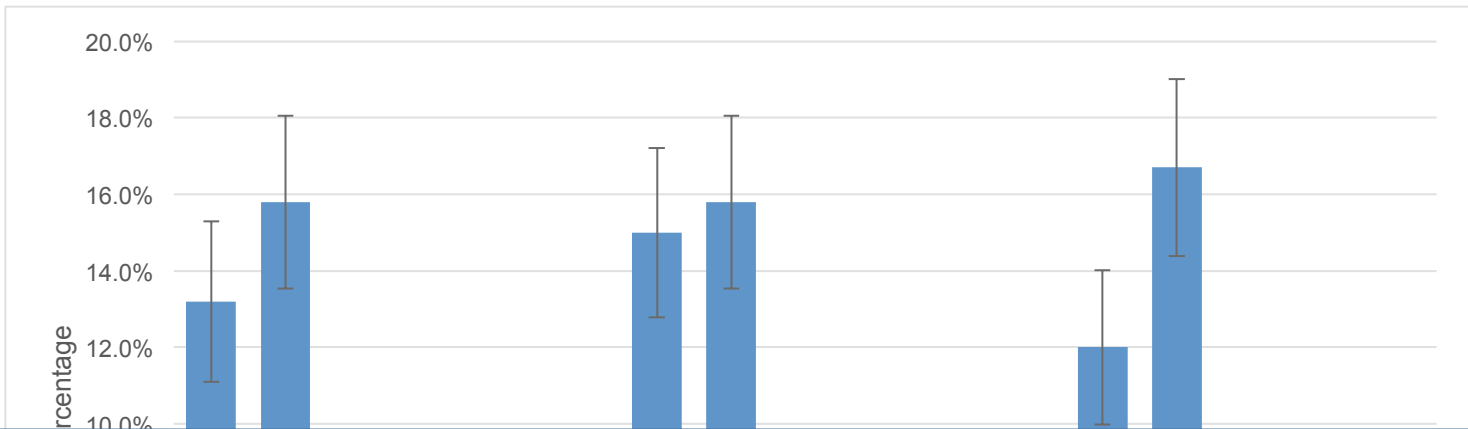
Fault Injection: Insn. Categories

Instruction category	LLFI selection criteria	PINFI selection criteria
<i>arithmetic</i>	Instructions that perform arithmetic or logical operations	Instructions that perform arithmetic or logical operations
<i>cast</i>	Instructions with 'cast' opcode	Instructions with 'convert' category
<i>cmp</i>	'cmp' instructions	Instructions whose next instruction is conditional branch
<i>load</i>	'load' instructions	'mov' instructions with memory as the source and register as the destination
<i>all</i>	All instructions	All instructions

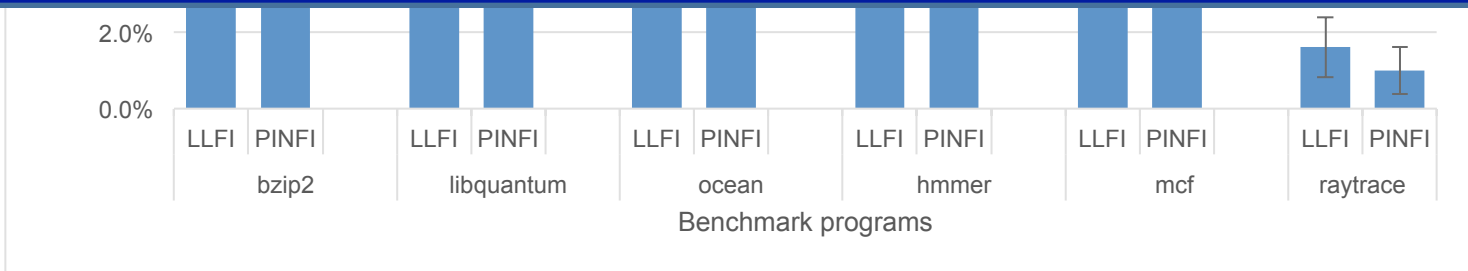
Results: Overall Failure Distribution



Results: SDCs for all instructions

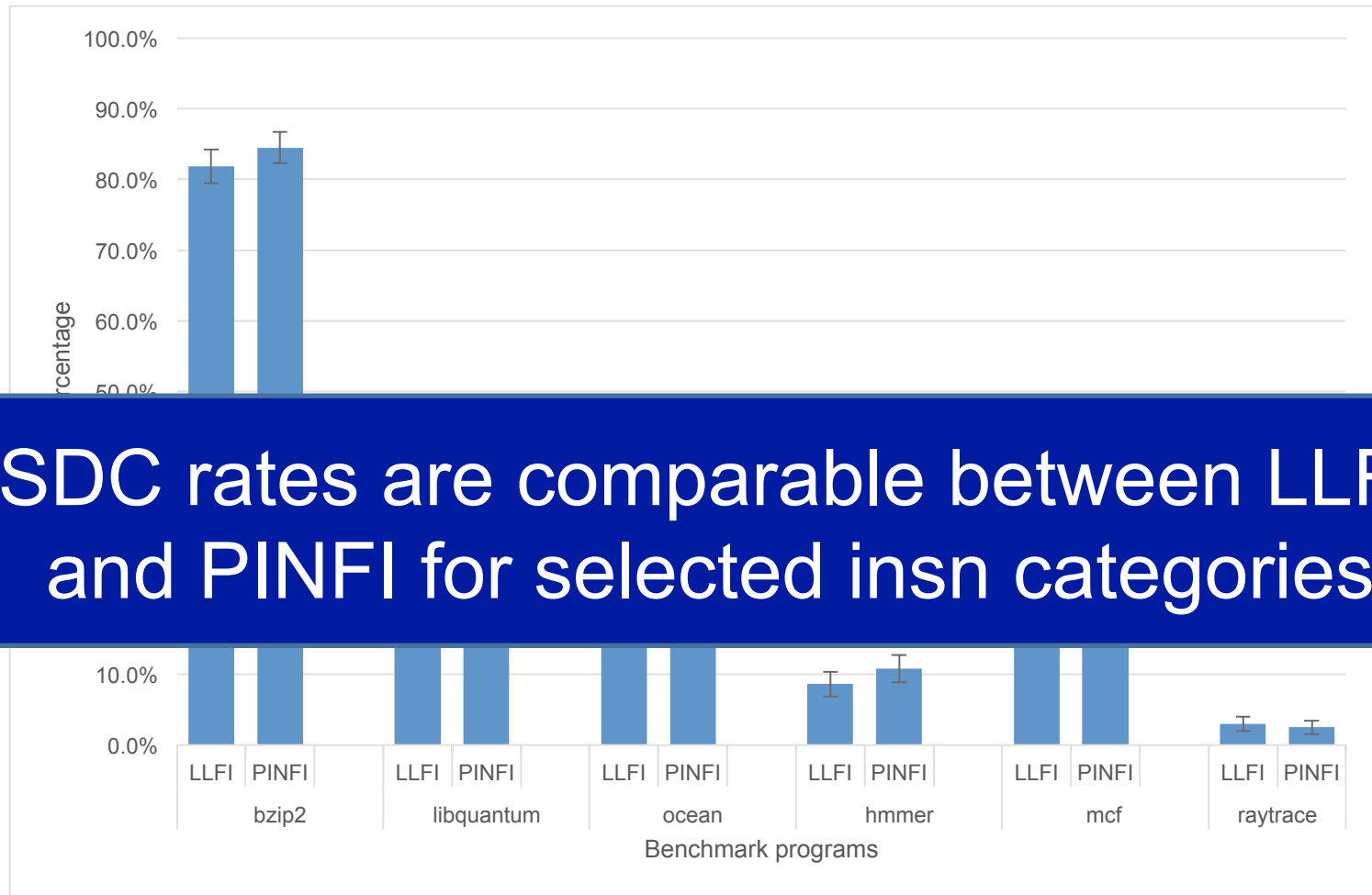


SDC rates are comparable between LLFI and PINFI for “all instructions”



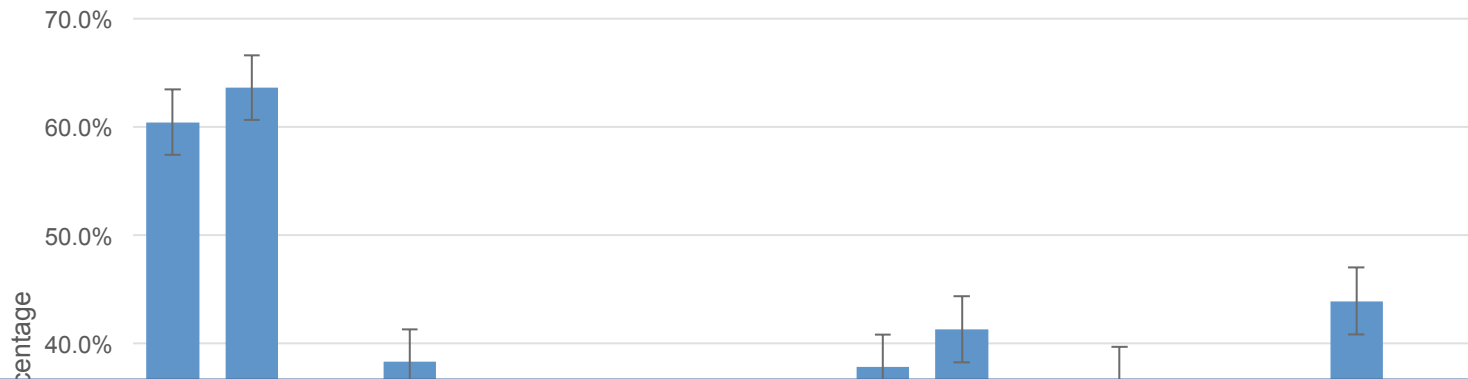
Error bars are computed at the 95% confidence level

Results: SDCs for 'cmp' instructions

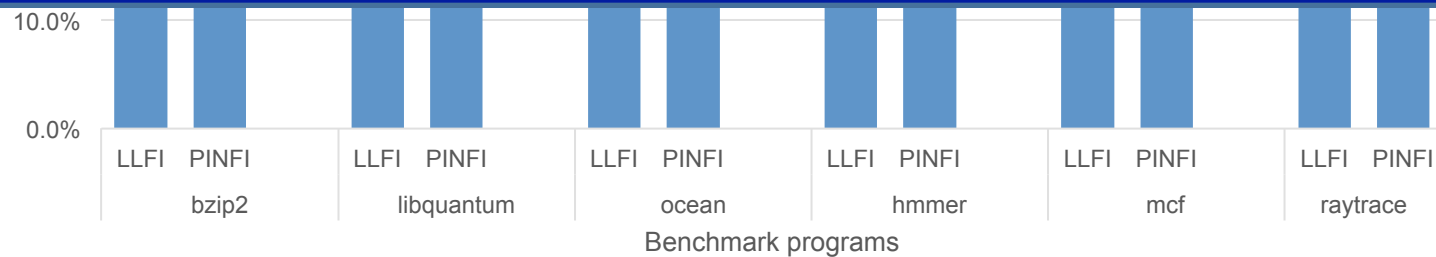


Error bars are computed at the 95% confidence level

Results: Crashes for all instructions



Crash rates differ widely between LLFI and PINFI for “all instructions”



Error bars are computed at the 95% confidence level

Why do crashes have poor accuracy in LLFI ?

- **Pointer computations in LLVM IR**
 - Abstracted away by GetElementPtr Instruction
 - Some pointer computations are a part of the instructions' encoding in assembly code
- **Mov instructions in x86 assembly code can move data between memory and registers**
 - Represented by loads and stores in LLVM IR
 - Some mov instructions are due to register spills

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Conclusion

- **Evaluate accuracy of high-level fault injection**
 - LLFI¹ as the high-level fault injector
 - PINFI² as the low-level fault injector
- **Results for accuracy of high-Level injection**
 - Accurate for SDC causing errors
 - Inaccurate for crash causing errors

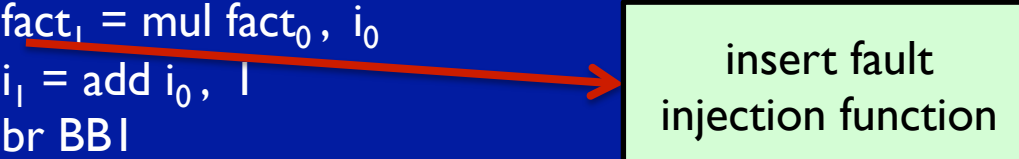
[1. https://github.com/DependableSystemsLab/LLFI](https://github.com/DependableSystemsLab/LLFI)

[2. https://github.com/DependableSystemsLab/PINFI](https://github.com/DependableSystemsLab/PINFI)

LLFI Framework: Operation

```
int main() {  
    int fact, i, n;  
    n = atoi (argv[1]);  
    fact = 1;  
  
    for( i = 1 ; i <= n; i++ )  
        fact = fact * i;  
  
    print fact;  
}
```

```
int main() {  
entry:  
    n1 = atoi (argv[1]);  
    br BBI  
  
BB:  
    fact1 = mul fact0, i0  
    i1 = add i0, 1  
    br BBI  
  
BBI:  
    i0 = phi [1, entry] , [i1, BB]  
    fact0 = phi [1, entry], [fact1, BB]  
    cond = sle i0, n1  
    br cond, label BB, label Return  
  
Return:  
    print fact0 }
```



insert fault
injection function

LLFI Framework: Operation

```
int main() {  
    int fact, i, n;  
    n = atoi (argv[1]);  
    fact = 1;  
  
    for( i = 1 ; i <= n; i++ )  
        fact = fact * i;  
  
    print fact;  
}
```

```
int main() {  
entry:  
    n1 = atoi (argv[1]);  
    br BBI  
  
BB:  
    fact1 = mul fact0, i0  
    fiI0 = call inject(I0, fact1)  
    i1 = add i0, 1  
    br BBI  
  
BBI:  
    i0 = phi [I, entry] , [i1, BB]  
    fact0 = phi [I, entry], [fiI0, BB]  
    cond = sle i0, n1  
    br cond, label BB, label Return  
  
Return:  
    print fact0 }
```

Replace all uses of
original with return val

